



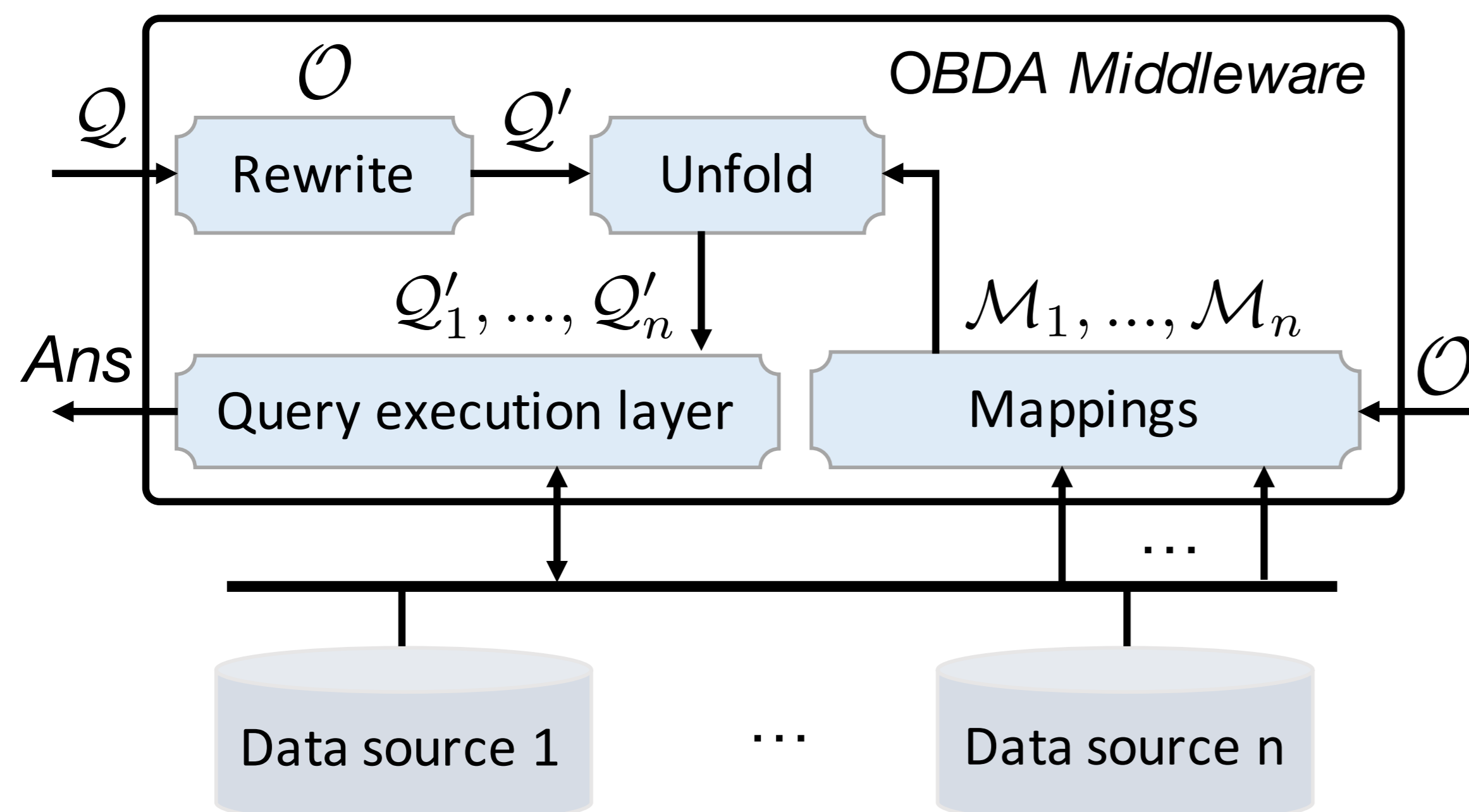
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Ontology-mediated Analytic Query Answering

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Analytics-aware ontology-based data access

OBDA – Classical setting



Limitations and challenges

- *Rewritability* of queries is not always possible
- Limited query language (e.g., no support for analytical operations like min, max, count, sum, avg)
- *Semantic mismatch* between data sources and ontological level

Data sources	Ontological level
closed world (CWA)	open world (OWA)
bag semantics	vs. set semantics
null values	existential quantification

- ⇒ Problematic with respect to supporting analytical operations
- *Optimisation* of ontological queries is *hard*

Approach

Motivating example

\mathcal{O} : $Sensor \sqsubseteq \exists hasMeasmnt$
 $Sensor \sqsubseteq \exists hasMUnit$

SELECT sid FROM Sensors \rightsquigarrow $Sensor(sid)$

\mathcal{M} : SELECT sid, unit FROM Sensors \rightsquigarrow $hasMUnit(sid, unit)$
SELECT sid, val FROM Measmnts \rightsquigarrow $hasMeasmnt(sid, val)$

Sensors (sid,tid,unit)			Measmnts (sid,val,t)		
sensor1	turbine1	celsius	sensor1	-50	t1
sensor2	turbine1	psi	sensor2	5000	t1
sensor3	turbine2	rpm	sensor3	2500	t1
...			sensor1	-50	t2
			...		

What is the answer to the following questions?

Q_1 : What is the number of measurements?

Q_2 : What is the maximum measurement value for each sensor?

Query evaluation

- ACQ query language: CQ + aggregate functions in head [1, 2]
 $Q_1 : q(\text{count}(y)) \leftarrow \exists x hasMeasmnt(x, y)$
 $Q_2 : q(x, \text{max}(y)) \leftarrow hasMeasmnt(x, y)$
- **Rewritability**: Are there Q'_1, Q'_2 over the data sources to answer Q_1 and Q_2 over \mathcal{O} ?

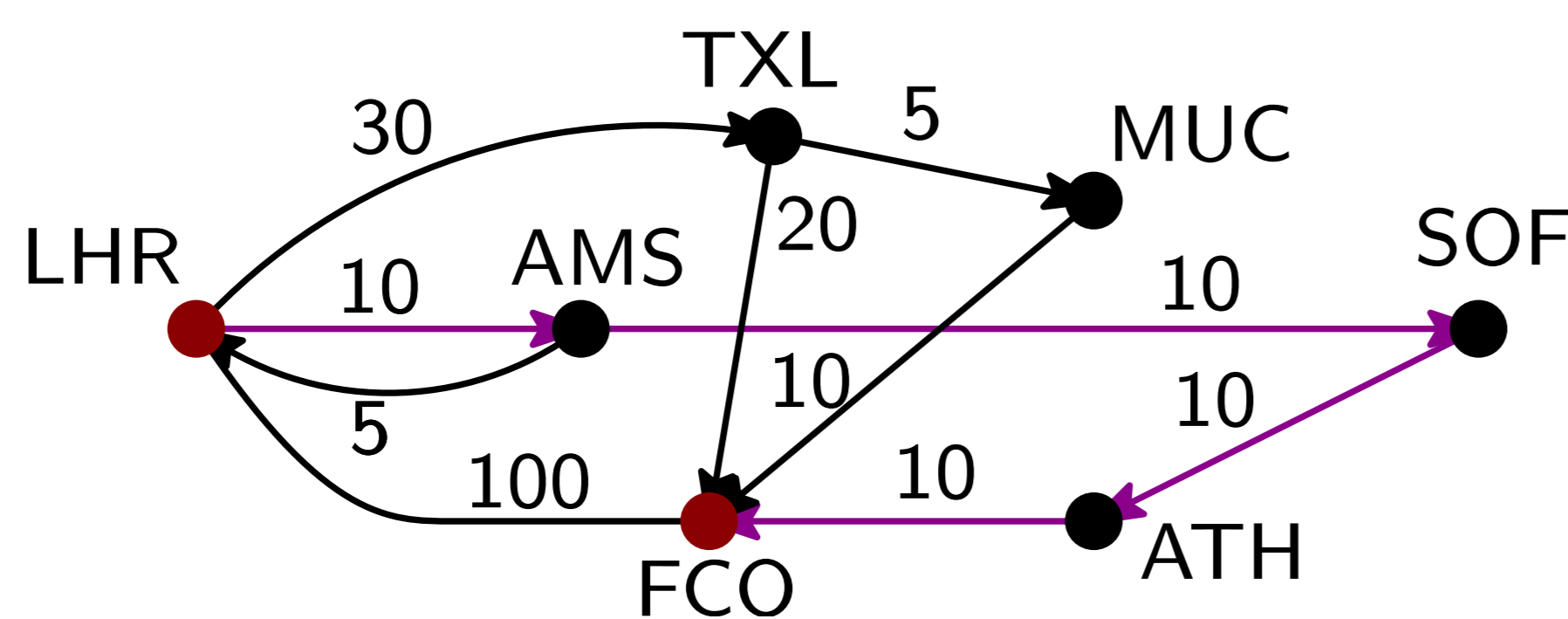
Semantics

- Extend DL-Lite_A with *multi-set* (bag) semantics
⇒ conservative extension of set semantics
⇒ preserves cardinality from data sources
- count-ACQs: (a) **minimal** [2] and (b) **skolem-based** semantics
⇒ identify *safe queries* for which minimality is tractable
- non-count-ACQs: aggregate over known values [1]

Expressiveness: pushing the envelope further

Can we answer the following questions using ACQs?

- What is the lowest cost for flying from London to Rome?
- How many bolts does a car manufacturer need to order for a specific model?



Challenges

- Apart from *aggregation* (e.g., sum, min, and count), we need also *recursion* to formulate/answer the above questions
- Assume a *rule-based language* extended with *aggregates*
- Interaction between recursion and aggregation is very powerful and non-trivial to manage

Previous work

- No semantics for the general case
- Proposals in the literature unsatisfactory:
* High complexity, no value invention [3]
* Undecidability of fact entailment [4, 5, 6]
* Limited expressivity (e.g., functionality) [3, 6, 7]
* Unnatural syntactic restrictions (hard to write programs) [4]

Our goals

- Define intuitive semantics leading to a *unique model*
- Offer a natural and *user-friendly syntax*
- *Generalise* existing approaches
- *Low complexity* of query evaluation; *sufficient* expressive power

Example program for cheapest flights

$\text{flight}(X, Z, C) \leftarrow \text{flight}(X, Y, C_1), \text{flight}(Y, Z, C_2), C = C_1 + C_2$
 $\text{cheapest flight}(X, Y, C) \leftarrow C = \min(C, \text{flight}(X, Y, C))$

References

- [1] D. Calvanese et al. Aggregate queries over ontologies. In *ONISW '08*.
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- [3] W. Faber et al. Semantics and complexity of recursive aggregates in answer set programming. *AI '11*.
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- [5] A. Van Gelder. Foundations of aggregation in deductive databases. In *DOOD'93*.
- [6] S. Greco. Dynamic programming in datalog with aggregates. *TKDE '99*.
- [7] I. S. Mumick et al. The magic of duplicates and aggregates. In *VLDB '90*.